

# Characterizations of MRD and Gabidulin codes

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In collaboration with Kyle Marshall.

## Introduction

- *Rank-distance codes* are matrix codes  $C \subseteq \mathbb{F}_q^{m \times n}$ , equipped with the rank distance

$$d_R(A, B) = \text{rank}(A - B).$$

They are useful in network coding, distributed storage ect.

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- Singleton bound for linear rank-distance codes  $C \subseteq \mathbb{F}_q^{m \times n}$  of dimension  $k$ :

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$$d_{\min}(C) \leq \min(m, n) - \max(m, n)k + 1.$$

- Codes that attain this bound exist for any set of parameters and are called *maximum rank distance (MRD) codes*.

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  - John Sheekey presented a construction related to semi-fields at BIRS workshop in January (and today).
  - The previous talk by Wolfgang Willems et al.
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- *Our contribution:* We will give some general characterizations of linear MRD and Gabidulin codes. The results give rise to equations one can solve to find all MRD or Gabidulin codes (up to equivalence).

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- We will assume  $n \leq m$  and study MRD codes  $C \subseteq \mathbb{F}_{q^m}^n$  that are  $\mathbb{F}_{q^m}$ -linear. These codes have a generator matrix  $G \in \mathbb{F}_{q^m}^{k \times n}$  and a respective parity check matrix  $H \in \mathbb{F}_{q^m}^{(n-k) \times n}$ .

### Example

Let  $\mathbb{F}_2^2 = \mathbb{F}_2[\alpha]$  and

$$G = (1, \alpha).$$

Then the code generated by  $G$  is

$$C = \{(0, 0), (1, \alpha), (\alpha, \alpha^2), (\alpha^2, 1)\}$$

$$\cong \begin{pmatrix} 0 & 0 \\ 0 & 0 \end{pmatrix}, \begin{pmatrix} 1 & 0 \\ 0 & 1 \end{pmatrix}, \begin{pmatrix} 0 & 1 \\ 1 & 1 \end{pmatrix}, \begin{pmatrix} 1 & 1 \\ 1 & 0 \end{pmatrix}.$$

It has dimension 1 (over  $\mathbb{F}_{2^2}$ ) and minimum rank distance 2 (over  $\mathbb{F}_2$ ). A respective parity check matrix is

$$H = (\alpha, 1).$$

## Known characterization of the MRD property

### Theorem (Gabidulin)

*Let  $H \in \mathbb{F}_{q^m}^{(n-k) \times n}$  be a parity check matrix of the code  $C$ . Then  $C$  is MRD if and only if*

$$\text{rank}(VH^T) = n - k$$

*for all  $V \in \mathbb{F}_q^{(n-k) \times n}$  with  $\text{rank}(V) = n - k$ .*

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*for all  $V \in \mathbb{F}_q^{(n-k) \times n}$  with  $\text{rank}(V) = n - k$ .*

*Simplification:* Since  $\text{GL}_{n-k}(q)$  does not change the rank of  $VH^T$ , it suffices to check the rank property for all elements of the left orbit of  $H^T$  under  $\mathcal{G}_q(n-k, n)$  (i.e. only  $V$  in reduced row echelon form).

## Towards a new characterization of the MRD property

It is known that  $\mathrm{GL}_n(q)$  acts (on the right) as an isometry on rank-metric codes in  $\mathbb{F}_{q^m}^n$ .  $\implies$

### Lemma

*The orbit under  $\mathrm{GL}_n(q)$  of an MRD code in  $\mathbb{F}_{q^m}^n$  consists of just MRD codes.*

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### Lemma

*The generator matrix  $G$  of an MRD code  $C \subseteq \mathbb{F}_{q^m}^n$  of dimension  $k$  in reduced row echelon form is of the form*

$$G = \left( I_k \mid A \right),$$

*where  $A \in (\mathbb{F}_{q^m} \setminus \mathbb{F}_q)^{k \times (n-k)}$ .*

### Theorem

*A generator matrix  $G \in \mathbb{F}_{q^m}^{k \times n}$  gives rise to an MRD code if and only if any element of the orbit of  $G$  under  $\text{GL}_n(q)$  has only non-zero maximal minors.*



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*Simplification:* Instead of all of  $\text{GL}_n(q)$  it suffices to study the orbit of the subgroup of

- the upper triangular matrices (since swapping columns does not change the minors, up to sign)
- with an all-1 diagonal (since multiplying columns of the generator matrix with  $\mathbb{F}_q^*$ -scalars does not change the non-zero property of the minors).

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## Definition

Let  $g_1, \dots, g_n \in \mathbb{F}_{q^m}$  be linearly independent over  $\mathbb{F}_q$ . The code with generator matrix

$$G = \begin{pmatrix} g_1 & g_2 & \dots & g_n \\ g_1^q & g_2^q & \dots & g_n^q \\ g_1^{q^2} & g_2^{q^2} & \dots & g_n^{q^2} \\ \vdots & \vdots & & \vdots \\ g_1^{q^{k-1}} & g_2^{q^{k-1}} & \dots & g_n^{q^{k-1}} \end{pmatrix}$$

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*Gabidulin codes are MRD codes.*

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*Generalized Gabidulin codes are MRD codes.*

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### Theorem (generalized)

*An MRD code is a generalized Gabidulin code with parameter  $s$  if and only if*

$$\dim(C \cap C^{q^s}) = k - 1.$$



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*The dual of an MRD code is an MRD code. Moreover, the dual of a Gabidulin code is a Gabidulin code.*

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*All MRD codes of dimension  $k = n - 1$  are Gabidulin codes.*

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*All MRD codes of length  $n = 1, 2, 3$  are Gabidulin codes.*

**First non-trivial case**  $n = m = 4, k = 2$ 

We want to find a description of all MRD codes with generator matrix in RREF

$$G = \begin{pmatrix} 1 & 0 & a & b \\ 0 & 1 & c & d \end{pmatrix}, \quad a, b, c, d \in \mathbb{F}_{q^4} \setminus \mathbb{F}_q.$$

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$$G = \begin{pmatrix} 1 & 0 & a & b \\ 0 & 1 & c & d \end{pmatrix}, \quad a, b, c, d \in \mathbb{F}_{q^4} \setminus \mathbb{F}_q.$$

We require that the orbit of upper-triangular matrices gives matrices with only non-zero maximal minors:

$$G \begin{pmatrix} 1 & u_1 & u_2 & u_3 \\ 0 & 1 & u_4 & u_5 \\ 0 & 0 & 1 & u_6 \\ 0 & 0 & 0 & 1 \end{pmatrix} = \begin{pmatrix} 1 & u_1 & u_2 + a & u_3 + au_6 + b \\ 0 & 1 & u_4 + c & u_5 + cu_6 + d \end{pmatrix}$$

for any  $u_1, \dots, u_6 \in \mathbb{F}_q$ .

Now we want to find a description of all non-Gabidulin MRD codes. For this we need that  $\dim(C \cap C^q) \neq k - 1$  (and  $a, b, c, d \in \mathbb{F}_{q^4} \setminus \mathbb{F}_q$ ):

$$\begin{aligned} & \text{rank} \begin{pmatrix} 1 & 0 & a & b \\ 0 & 1 & c & d \\ 1 & 0 & a^q & b^q \\ 0 & 1 & c^q & d^q \end{pmatrix} \neq 3 \\ & \iff \text{rank} \begin{pmatrix} 1 & 0 & a & b \\ 0 & 1 & c & d \\ 0 & 0 & a^q - a & b^q - b \\ 0 & 0 & c^q - c & d^q - d \end{pmatrix} = 4 \\ & \iff (a^q - a)(d^q - d) - (b^q - b)(c^q - c) \neq 0. \end{aligned}$$



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For non-generalized Gabidulin codes with  $s = 3$  we get

$$(a^{q^3} - a)(d^{q^3} - d) - (b^{q^3} - b)(c^{q^3} - c) \neq 0.$$

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### Theorem

*All  $2^4$ -linear MRD codes in  $\mathbb{F}_{2^4}^4$  are Gabidulin codes.*

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### Theorem

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For  $q = 3$  we found many solutions, e.g. (with  $\alpha^4 = \alpha^3 + 1$ )

$$G = \begin{pmatrix} 1 & 0 & \alpha & \alpha^2 \\ 0 & 1 & \alpha^2 & 2\alpha \end{pmatrix}.$$

### Theorem

*There are many non-Gabidulin  $3^4$ -linear MRD codes in  $\mathbb{F}_{3^4}^4$ .*

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- With these criteria we show that all  $q^m$ -linear MRD codes in  $\mathbb{F}_{q^m}^n$  are (generalized) Gabidulin codes if
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- Moreover, we give an example of a non-Gabidulin  $q^m$ -linear MRD code for  $q = 3, n = m = 4, k = 2$ .



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- Moreover, we give an example of a non-Gabidulin  $q^m$ -linear MRD code for  $q = 3, n = m = 4, k = 2$ .
- This method can also be used for slightly larger parameters.

Thank you for your attention!



Questions?   Comments?