

NEW RESULTS ON GRIESMER CODES AND ARCS

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1. Linear Codes over Finite Fields

- ◇ **Linear $[n, k]_q$ code**: $C < \mathbb{F}_q^n$, $\dim C = k$
- ◇ **$[n, k, d]_q$ -code**: $d = \min\{d(\mathbf{u}, \mathbf{v}) \mid \mathbf{u}, \mathbf{v} \in C, \mathbf{u} \neq \mathbf{v}\}$.
 - n - the **length** of C ;
 - k - the **dimension** of C ;
 - d - the **minimum distance** of C .
- ◇ A_i – number of codewords of (Hamming) weight i
- ◇ $(A_i)_{i \geq 0}$ – the **spectrum** of C

The Main Problem in Coding Theory.

Optimize one of the parameters n , k , d , given the other two.

$n_q(k, d)$ - minimal length of a linear code over \mathbb{F}_q of dimension k and minimum distance d ;

$K_q(n, d)$ - maximal dimension of a linear code over \mathbb{F}_q of length n and minimum distance d ;

$D_q(n, k)$ - maximal minimum distance of a linear code over \mathbb{F}_q of length n and dimension k .

optimality with respect to $n \implies$ optimality with respect to k and d

◇ Griesmer bound: Let \mathcal{C} be an $[n, k, d]_q$ -code. Then

$$n_q(k, d) \geq g_q(k, d) = \sum_{i=0}^{k-1} \left\lceil \frac{d}{q^i} \right\rceil$$

Theorem. Given the integer k and the prime power q , Griesmer $[g_q(k, d), k, d]_q$ -codes exist for all sufficiently large d .

The problem of finding the exact value of $n_q(k, d)$ is solved for

- $q = 2$: $k \leq 8$ for all d ;
- $q = 3$: $k \leq 5$ for all d ;
- $q = 4$: $k \leq 4$ for all d ;
- $q = 5, 7, 8, 9$: $k \leq 3$ for all d ;
- $q = 5$: $k = 4$ – four values of d for which $n_5(4, d)$ is not known.

<http://www.mi.s.osakafu-u.ac.jp/maruta/griesmer.htm>

The Open Cases for $q = 5, k = 4$

| d | $g_5(4, d)$ | $n_5(4, d)$ | \mathcal{K} | $\mathcal{K} _H$ |
|-----|-------------|-------------|---------------|------------------|
| 81 | 103 | 103–104 | (103, 22) | (22, 5)-arc |
| 82 | 104 | 104–105 | (104, 22) | in PG(2, 5) |
| 161 | 203 | 203–204 | (203, 42) | (42, 9)-arc |
| 162 | 204 | 204–205 | (204, 42) | in PG(2, 5) |

2. Divisible and Quasidivisible Arcs

◇ A **multiset** in $\text{PG}(k-1, q)$ is a mapping

$$\mathcal{K} : \begin{cases} \mathcal{P} & \rightarrow \mathbb{N}_0, \\ P & \rightarrow \mathcal{K}(P). \end{cases}$$

◇ $\mathcal{K}(P)$ – **multiplicity** of the point P .

◇ $\mathcal{Q} \subset \mathcal{P}$: $\mathcal{K}(\mathcal{Q}) = \sum_{P \in \mathcal{Q}} \mathcal{K}(P)$ – **multiplicity** of the set \mathcal{Q} .

◇ $\mathcal{K}(\mathcal{P})$ – the **cardinality** of \mathcal{K} .

◇ Points, lines, ... , hyperplanes of multiplicity i are called i -points, i -lines, ... , i -hyperplanes.

◇ a_i – the number of hyperplanes H with $\mathcal{K}(H) = i$

◇ $(a_i)_{i \geq 0}$ – the **spectrum** of \mathcal{K}

Definition. (n, w) -arc in $\text{PG}(k - 1, q)$: a multiset \mathcal{K} with

- 1) $\mathcal{K}(\mathcal{P}) = n$;
- 2) for every hyperplane H : $\mathcal{K}(H) \leq w$;
- 3) there exists a hyperplane H_0 : $\mathcal{K}(H_0) = w$.

Definition. (n, w) -blocking set in $\text{PG}(k - 1, q)$

(or (n, w) -minihyper): a multiset \mathcal{K} with

- 1) $\mathcal{K}(\mathcal{P}) = n$;
- 2) for every hyperplane H : $\mathcal{K}(H) \geq w$;
- 3) there exists a hyperplane H_0 : $\mathcal{K}(H_0) = w$.

Definition. An (n, w) -arc \mathcal{K} in $\text{PG}(k - 1, q)$ is called t -extendable, if there exists an $(n + t, w)$ -arc \mathcal{K}' in $\text{PG}(k - 1, q)$ with $\mathcal{K}'(P) \geq \mathcal{K}(P)$ for every point $P \in \mathcal{P}$. An 1-extendable arc is called extendable.

Definition. An arc \mathcal{K} in $\text{PG}(k - 1, q)$ with $\mathcal{K}(\mathcal{P}) = n$ and spectrum (a_i) is said to be divisible with divisor Δ , $\Delta > 1$, if $a_i = 0$ for all $i \not\equiv n \pmod{\Delta}$.

Definition. An arc \mathcal{K} with $\mathcal{K}(\mathcal{P}) = n$ and spectrum (a_i) is said to be t -quasidivisible with divisor Δ , $\Delta > 1$, (or t -quasidivisible modulo Δ) if $a_i = 0$ for all $i \not\equiv n, n + 1, \dots, n + t \pmod{\Delta}$.

3. Linear codes as multisets of points

| | | |
|---|-------------------|--|
| $[n, k, d]_q$ -code C of full length | \Leftrightarrow | $(n, w = n - d)$ -arc \mathcal{K} in $\text{PG}(k - 1, q)$ |
| $\mathbf{0} \neq \mathbf{u} \in C, \text{wt}(\mathbf{u}) = u$ | \Leftrightarrow | a hyperplane H with $\mathcal{K}(H) = n - u,$ |
| extendable $[n, k, d]_q$ -code C | \Leftrightarrow | extendable $(n, n - d)$ -arc \mathcal{K} |
| divisible $[n, k, d]_q$ -code $A_i = 0$ for all $i \not\equiv 0 \pmod{\Delta}$ | \Leftrightarrow | divisible $(n, n - d)$ -arc in $\text{PG}(k - 1, q)$ $a_i = 0$ for all $i \not\equiv n \pmod{\Delta}$ |
| t -quasidivisible $[n, k, d]_q$ -code $A_i = 0$ for all $i \not\equiv -j \pmod{q}$ $j \in \{0, 1, \dots, t\}$ | \Leftrightarrow | t -quasidivisible $(n, n - d)$ -arc in $\text{PG}(k - 1, q)$ $a_i = 0$ for all $i \not\equiv n + j \pmod{q}$ |

◇ Griesmer arcs: arcs associated with codes meeting the Griesmer bound

Griesmer $[n, k, d]_q$ codes \Leftrightarrow Griesmer (n, w) -arcs in $\text{PG}(k-1, q)$

$$n = \sum_{i=0}^{k-1} \lceil d/q^i \rceil$$

$$n = \sum_{i=0}^{k-1} \lceil (n-w)/q^i \rceil$$

◇ If $d = n - w = sq^{k-1} - \varepsilon_{k-2}q^{k-2} - \dots - \varepsilon_1q - \varepsilon_0$, and

$w_i :=$ maximal multiplicity of a subspace of codimension i , $i = 0, \dots, k-1$.

Then

$$w_i = sv_{k-i} - \varepsilon_{k-2}v_{k-i-1} - \dots - \varepsilon_{i+1}v_2 - \varepsilon_iv_1,$$

where $v_k = (q^k - 1)/(q - 1)$.

4. Some Extension Results

Theorem. (R. Hill, P. Lizak, 1995, geometric version) Let \mathcal{K} be a (n, w) -arc in $\text{PG}(k-1, q)$ with $\gcd(n-w, q) = 1$. Let further $\mathcal{K}(H) \equiv n$ or $w \pmod{q}$ for all hyperplanes H . Then \mathcal{K} is extendable to a divisible $(n+1, w)$ -arc in $\text{PG}(k-1, q)$. In particular, every 1-quasidivisible arc with divisor q is extendable.

Theorem. (T. Maruta, 2004, geometric version) Let \mathcal{K} be a 2-quasidivisible (n, w) -arc in $\text{PG}(k-1, q)$, $q \geq 5$, odd, with divisor q . Then \mathcal{K} is extendable to an $(n+1, w)$ -arc in $\text{PG}(k-1, q)$.

◇ \mathcal{K} - t -quasidivisible (n, w) -arc in $\Sigma = \text{PG}(k-1, q)$, i.e. for every hyperplane H , we have $\mathcal{K}(H) \equiv n, n+1, \dots, n+t \pmod{q}$, where $0 < t < q$ is an integer constant.

◇ Define an arc $\tilde{\mathcal{K}}$ in the dual space $\tilde{\Sigma}$

$$\tilde{\mathcal{K}} : \begin{cases} \mathcal{H} & \rightarrow \mathbb{N}_0, \\ H & \rightarrow \tilde{\mathcal{K}}(H) := n + t - \mathcal{K}(H) \pmod{q}. \end{cases}$$

where \mathcal{H} is the set of all hyperplanes of Σ .

Theorem. Let \mathcal{K} be an (n, w) -arc in $\Sigma = \text{PG}(k-1, q)$ which is t -quasidivisible modulo q , $t < q$.

Let

$$\tilde{\mathcal{K}} = \sum_{i=1}^c \chi_{\tilde{H}_i} + \tilde{\mathcal{K}'}$$

for some arc $\tilde{\mathcal{K}'}$ and c not necessarily different hyperplanes $\tilde{H}_1, \dots, \tilde{H}_c$ then \mathcal{K} is c -extendable. In particular, if $\tilde{\mathcal{K}}$ contains a hyperplane in its support then \mathcal{K} is extendable.

Theorem. Let \tilde{S} be a subspace of $\tilde{\Sigma}$ of positive dimension. Then $\tilde{\mathcal{K}}(\tilde{S}) \equiv t \pmod{q}$.

Theorem. (Landjev, Rousseva, Storme, 2014)

Let \mathcal{K} be a t -quasidivisible Griesmer arc in $\text{PG}(k-1, q)$ with parameters (n, w) , where

$$d = n - w = sq^{k-1} - \varepsilon_{k-2}q^{k-2} - \dots - \varepsilon_1q - \varepsilon_0.$$

Let further $\varepsilon_0 = t, \dots, \varepsilon_{k-2} < \sqrt{q}$. Then \mathcal{K} is t -extendable.

- ◇ $\tilde{\mathcal{K}}$ is a (tv_{k-1}, tv_{k-2}) -arc, where $v_k = \frac{q^k - 1}{q - 1}$
- ◇ $\tilde{\mathcal{K}}$ is a sum of t hyperplanes

5. $(t \bmod q)$ -Arcs

Definition. An arc \mathcal{F} is called a $(t \bmod q)$ -arc if

- all points have multiplicity $\leq t$;
- all subspaces S of positive dimension have multiplicity $\mathcal{F}(S) \equiv t \pmod{q}$.

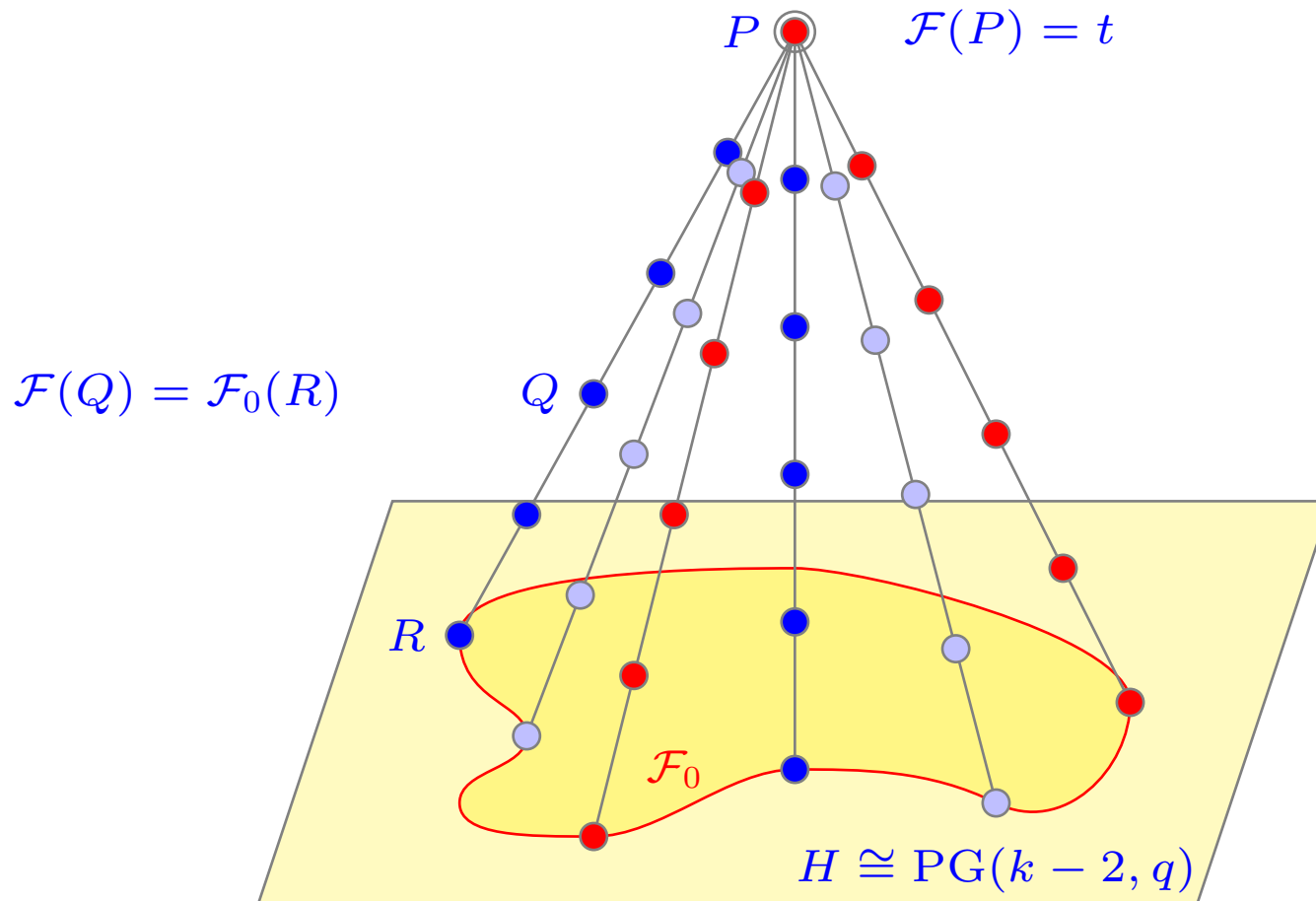
Theorem A. The sum of a $(t_1 \bmod q)$ -arcs and a $(t_2 \bmod q)$ -arc is a $(t \bmod q)$ -arc with $t = t_1 + t_2$. In particular, the sum of t hyperplanes in $\text{PG}(k-1, q)$ is a $(t \bmod q)$ -arc.

Theorem B. Let \mathcal{F}_0 be a $(t \bmod q)$ -arc in a hyperplane $H \cong \text{PG}(k-2, q)$ of $\Sigma = \text{PG}(k-1, q)$. For a fixed point $P \in \Sigma \setminus H$, define an arc \mathcal{F} in Σ as follows:

- $\mathcal{F}(P) = t$;
- for each point $Q \neq P$: $\mathcal{F}(Q) = \mathcal{F}_0(R)$ where $R = \langle P, Q \rangle \cap H$.

Then the arc \mathcal{F} is a $(t \bmod q)$ -arc in $\text{PG}(k-1, q)$ of size $q|\mathcal{F}_0| + t$.

Definition. $(t \bmod q)$ -arcs obtained by Theorem B are called *lifted arcs*.



\mathcal{F} : an arc in $\Sigma = \text{PG}(k-1, q)$

\mathcal{H} – the set of all hyperplanes in Σ

σ - a function such that $\sigma(\mathcal{F}(H))$ is a non-negative integer for all $H \in \mathcal{H}$.

The arc \mathcal{F}^σ in $\tilde{\Sigma}$

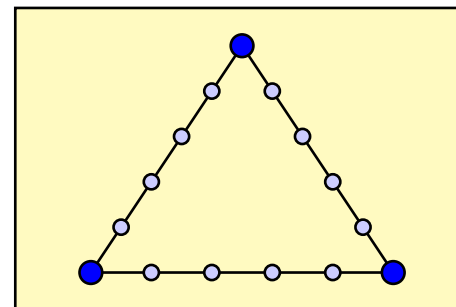
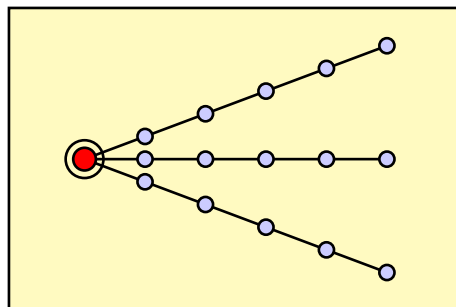
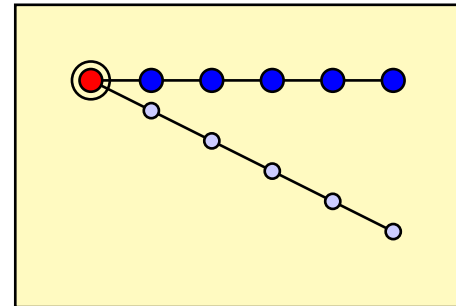
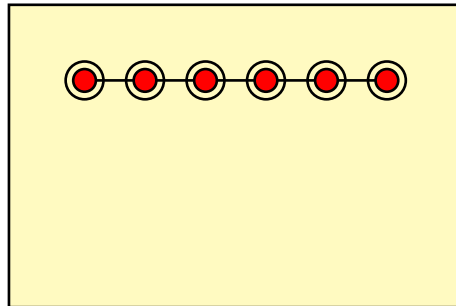
$$\mathcal{F}^\sigma : \begin{cases} \mathcal{H} & \rightarrow \mathbb{N}_0 \\ H & \rightarrow \sigma(\mathcal{F}(H)) \end{cases}$$

is called the σ -dual of \mathcal{F} .

Theorem C. Let \mathcal{F} be a $(t \bmod q)$ -arc in $\text{PG}(2, q)$ of size $mq + t$. Then the arc \mathcal{F}^σ with $\sigma(x) = (x - t)/q$ is an $((m - t)q + m, m - t)$ -blocking set in the dual plane with line multiplicities $m - t, m - t + 1, \dots, m$.

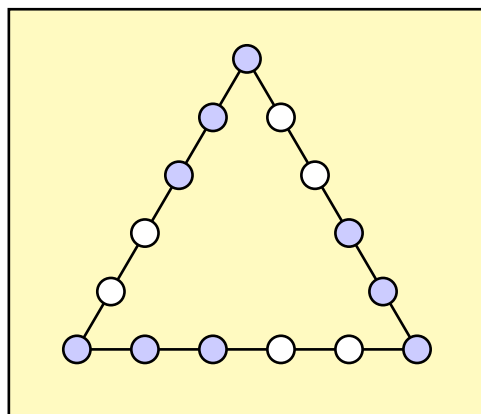
$(3 \pmod 5)$ -arcs in $PG(2, 5)$

$(18, \{3, 8, 13, 18\})$ -arcs



\mathcal{F} : $(23, \{3, 8\})$ -arc

\mathcal{F}^σ :



\mathcal{F}^σ : $(9,1)$ -blocking set

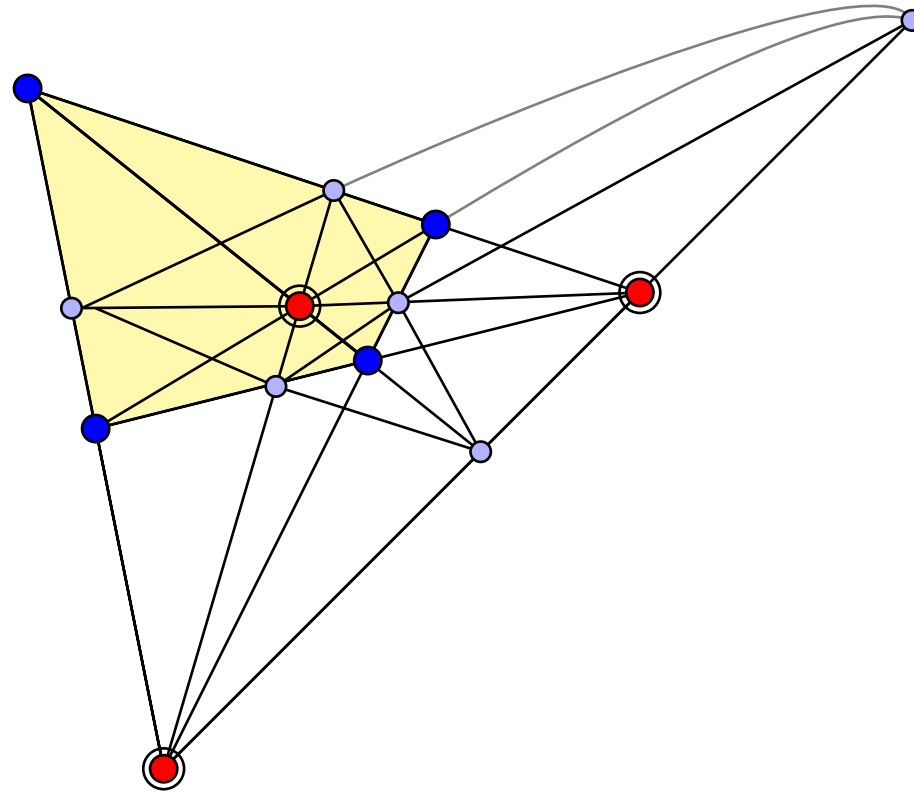
with line multiplicities 1, 2, 3, 4

\mathcal{F} : $(28, \{3, 8\})$ -arc

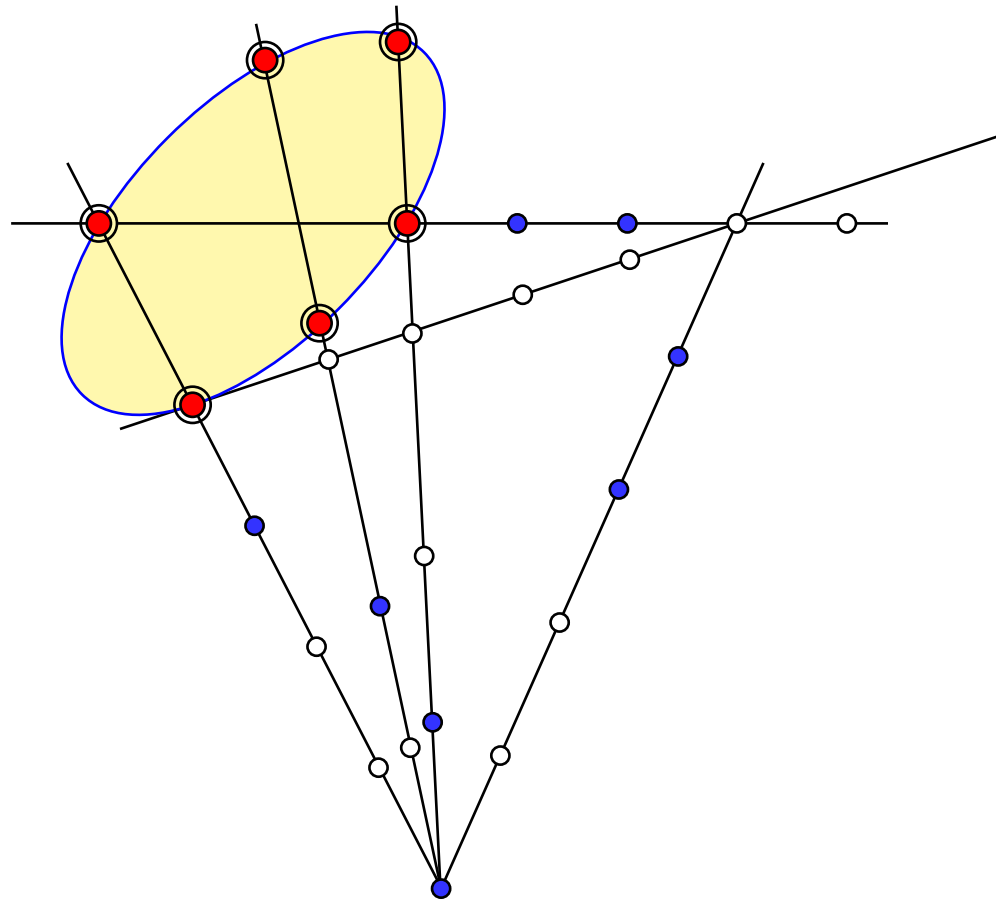
\mathcal{F}^σ : $(15, 2)$ -blocking set with line multiplicities 2, 3, 4, 5

\mathcal{F}^σ : the complement of the unique $(16, 4)$ -arc without external lines

The $(23, \{3, 8\})$ -arc



The $(28, \{3, 8\})$ -arc



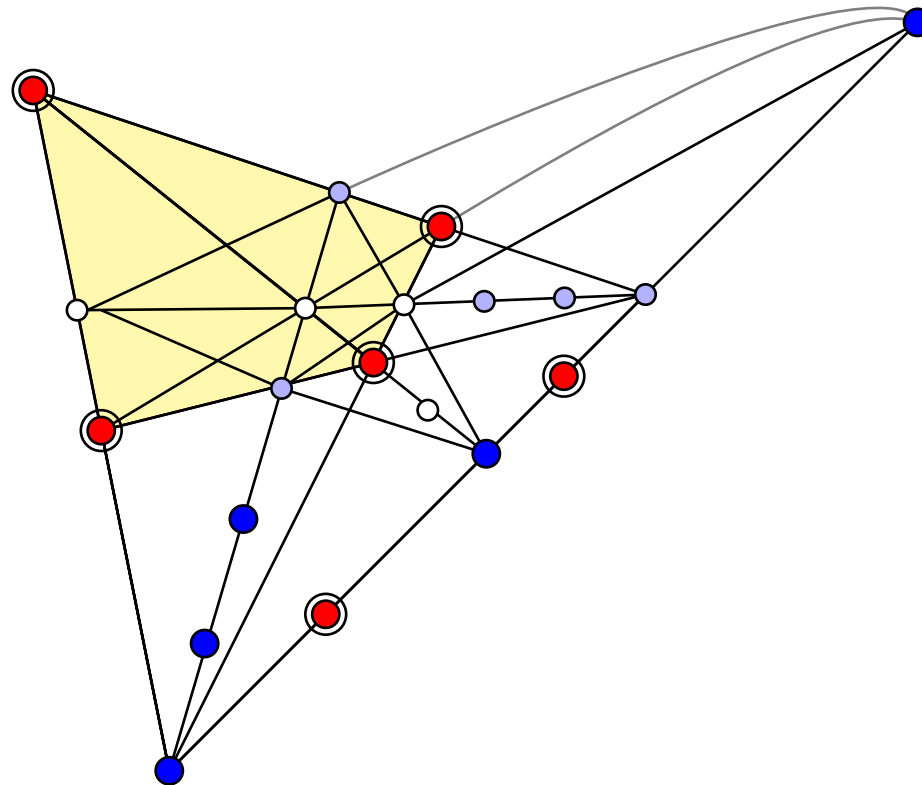
\mathcal{F} : (33, {3, 8, 13})-arc

\mathcal{F}^σ : (21, 3)-blocking set with line multiplicities 3, 4, 5, 6

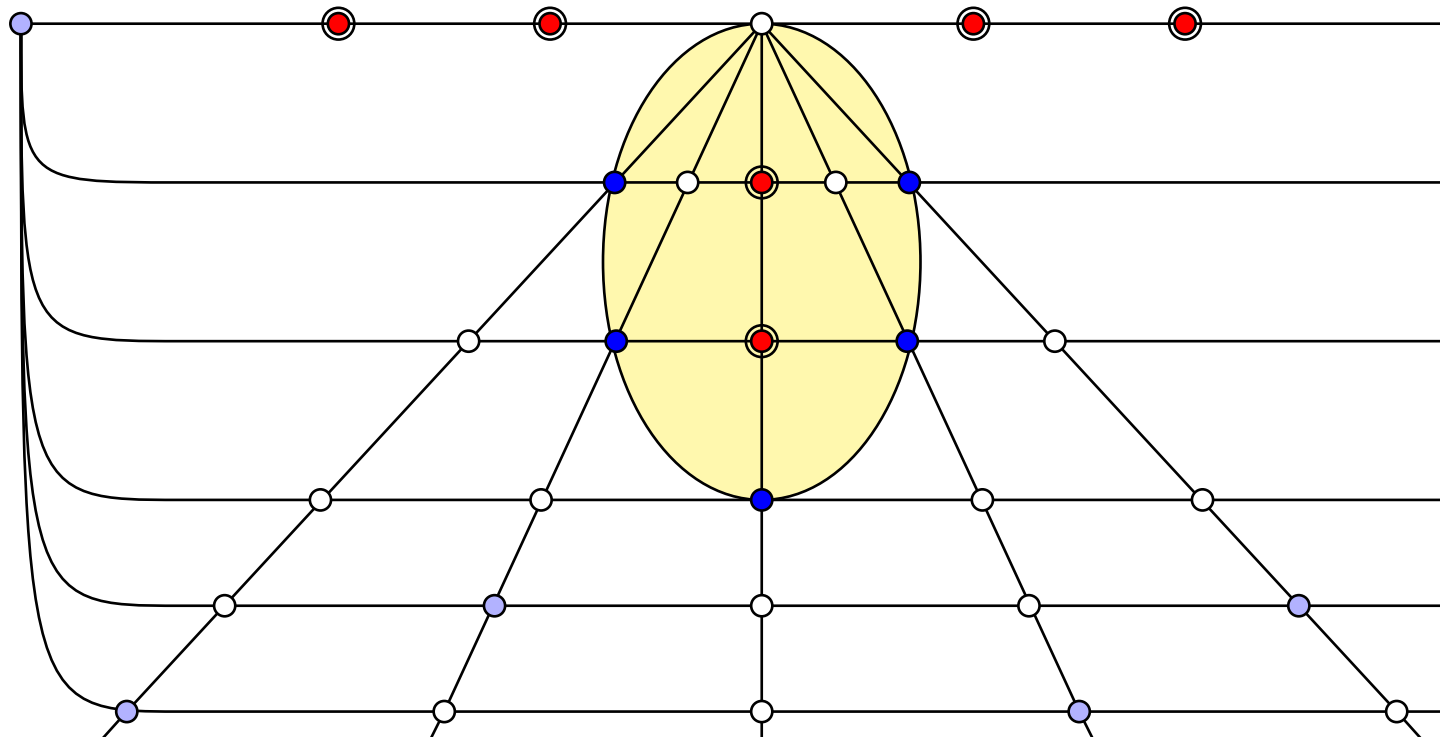
\mathcal{F}^σ is one of the following:

- (1) the complement of the seven non-isomorphic (10, 3)-arcs; $\Lambda_2 = 0$
- (2) the complement of the (11, 3)-arc with four external lines; a point not on an external line is doubled; $\Lambda_2 = 1$
- (3) one double point which forms an oval with five of the 0-points; the tangent in the 2-point is a 3-line; $\Lambda_2 = 1$
- (4) $\text{PG}(2, 5)$ minus a triangle with vertices of multiplicity 2, 2, 1; $\Lambda_2 = 2$

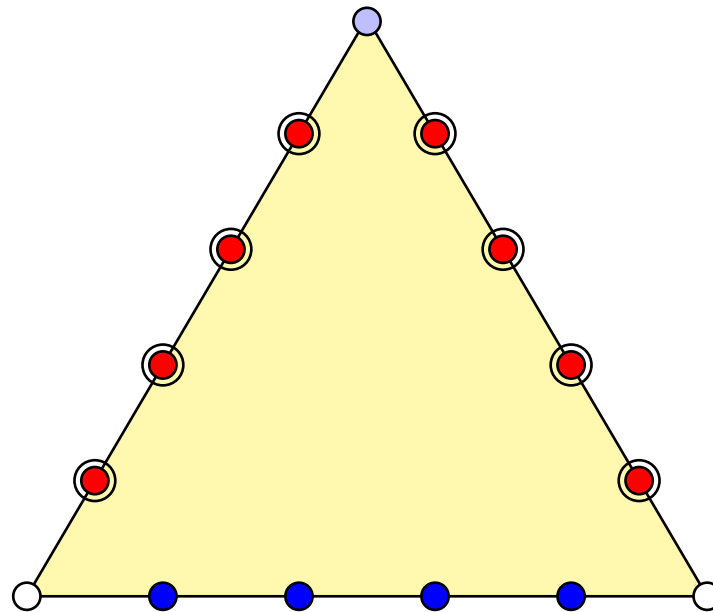
(2) The first $(33, \{3, 8, 13\})$ -arc with one 13-line



(3) the second $(33, \{3, 8, 13\})$ -arc with one 13-line



(4) $(33, \{3, 8, 13\})$ -arc with two 13-lines



Theorem D. Every $(3 \pmod 5)$ -arc \mathcal{F} in $\text{PG}(3, 5)$ with $|\mathcal{F}| \leq 153$ is a lifted arc (obtained by Theorem B). In particular, $|\mathcal{F}| = 93, 118$, or 143 .

6. The Nonexistence of Some Griesmer Codes

A. The Nonexistence of $[104, 4, 82]_5$ -Codes

$$d = 82 = 5^3 - 5^2 - 3 \cdot 5 - 3$$

$$s = 1, \varepsilon_2 = 1, \varepsilon_1 = 3, \varepsilon_0 = t = 3;$$

$$\varepsilon_0, \varepsilon_1 \geq \sqrt{q}$$

$$w_3 = 1, w_2 = 5, w_1 = 22, w_0 = n = 104$$

Theorem E. Let \mathcal{K} be a $(104, 22)$ -arc in $\text{PG}(3, 5)$. Then

- (i) \mathcal{K} is a Griesmer 3-quasidivisible projective arc;
- (ii) $\tilde{\mathcal{K}}$ is a $(3 \bmod 5)$ -arc in $\text{PG}(3, 5)$;
- (iii) there is no 18-plane π such that $\tilde{\mathcal{K}}|_{\pi}$ is the sum of three copies of the same line;
- (iv) $|\tilde{\mathcal{K}}| \leq 143$.

- Hence $|\tilde{\mathcal{K}}| = 93$ and $\tilde{\mathcal{K}}$ is a sum of three planes.
- Hence \mathcal{K} is 3-extendable to a (non-existent) $(107, 22)$ -arc.
- There is no $[104, 4, 82]_5$ -code and

$$n_5(4, 82) = 105.$$

B. The Nonexistence of $[q^3 - 3q - 6, 4, q^3 - q^2 - 3q - 3]_q$ -Codes

◇ $q = 5$: (104, 22)-arcs in $\text{PG}(3, 5)$

◇ $q = 7$: (316, 46)-arcs in $\text{PG}(3, 7)$

Plane multiplicities: 1, 8, 15, 22, 29, 36, 43–49

◇ $q = 8$: (482, 61)-arcs in $\text{PG}(3, 8)$

Plane multiplicities: 10, 42, 58–61

◇ $q = 9$: (696, 78)-arcs in $\text{PG}(3, 9)$

Plane multiplicities: 75–78

◇ $q \geq 11$: are 3-extendable since $\varepsilon_i < \sqrt{q}$.

C. The Nonexistence of $[204, 4, 162]_5$ -Codes